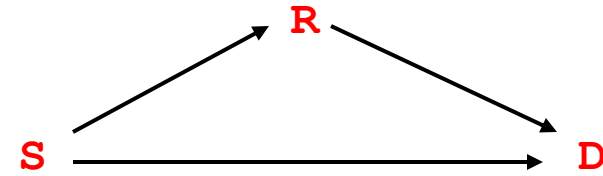


Practical Quantize-and-Forward schemes for the Frequency Division Relay Channel



PROBLEM STATEMENT

- Channel description:



- Information theoretic relay channel (RC)

- [Cover & El Gamal 1979]: Introduction of Estimate-and-Forward (EF) and Decode-and-Forward (DF) relaying schemes.
- [El Gamal & Al. 2006]: The optimal relaying scheme depends on channels SNRs.

- Relaying protocols:

- Decode-and-Forward

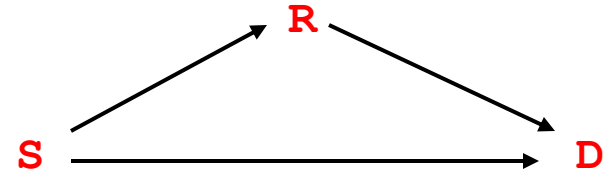
Forward decoding errors when the S-R link is bad.

- Amplify-and-Forward

$\alpha[hx + n] \rightarrow$ noise amplification

- Estimate-and-Forward

PROBLEM STATEMENT



○ Main issues

- Propose two relaying schemes that take SNRs into account:
 - Quantize-and-Forward (QF). Low complexity → scalar
 - Amplify-and-Forward with clipping (C-AF) .
- Performance analysis for different SNR scenarios.

○ Assumptions:

- Memoryless fading channels.
- Orthogonality between the downlink and the cooperative link.
- Coherent reception at R and D.
- CSI at R and D.
- SNRs knowledge at R and D.

BACKGROUND

- EF-based protocols for the RC
 - [Liu & Al. 2005] and [Charkrabarti & Al. 2006]: Rates analyses of proposed schemes for the non-orthogonal half-duplex RC.
 - [Hu & Al. 2006]: Practical EF schemes based on Wyner-Ziv coding for the orthogonal half-duplex RC.
- AF-based protocols for the RC
 - **Simplicity** → Develop advanced coding strategies (ex. Yang & Belfiore): DMT achieving.
 - No loss in terms on information but [analog constraint + **noise amplification**].
 - Optimal non-linear function at the relay in a multihop relay channel:
 - to minimize the BER [Abou-Faycal & Médard 2004]
→ **Complexity for high order modulation**
 - To maximize the mutual information [Gomadam & Jafar 2006]

OUTLINE

1. Signal model
2. Quantize-and-Forward
3. Maximum likelihood for the QF protocol
4. Clipped Amplify-and-Forward
5. Simulation results

SIGNAL MODEL

- **Source and relay: unit-power constraint**

$$E|X|^2 = 1 \quad E|X_r|^2 = 1$$

- **Block fading channels.**

- **Relay: received signal** $x_{sr}(n) = h_{sr} \times x(n) + w_{sr}(n)$

- **Destination: orthogonality leads to**

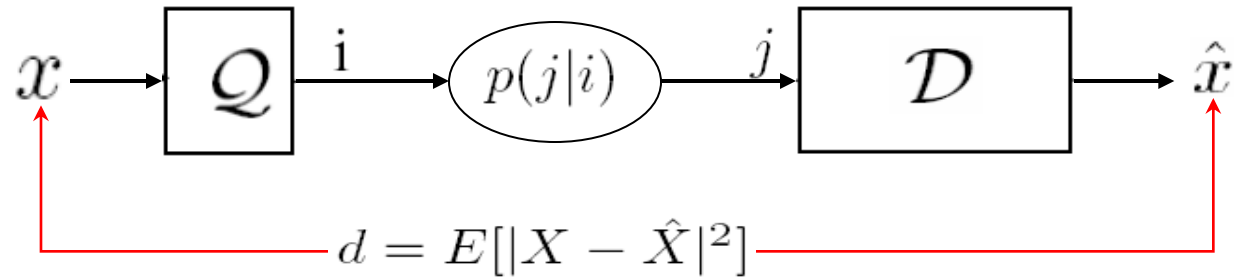
$$\begin{cases} y_{sd}(n) = h_{sd} \times x(n) + w_{sd}(n) \\ y_{rd}(n) = h_{rd} \times x_r(n) + w_{rd}(n) \end{cases}$$

Depends on the relaying scheme



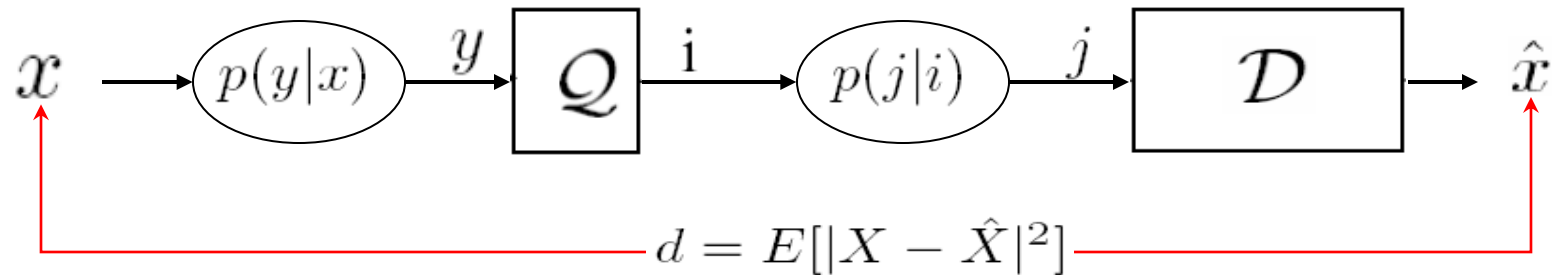
OPTIMIZED QF PROTOCOL (1/5)

○ Kunterbach and Wintz



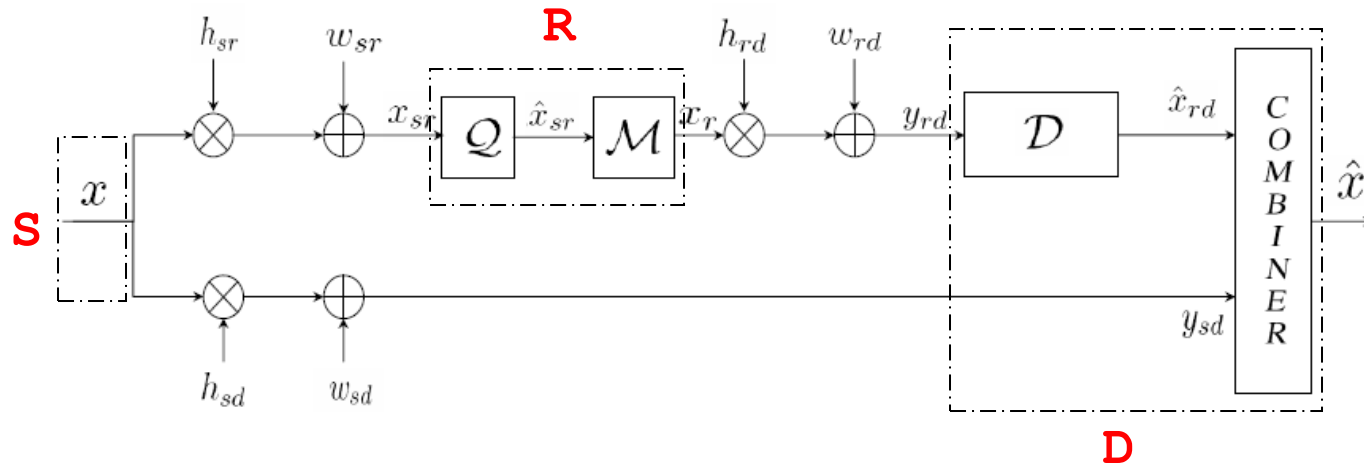
→ Modify the Lloyd-Max algorithm to take into account channel errors in the quantization optimization

○ RC: two-hop link S-R-D



→ Take into account S-R and R-D channels errors in the quantizer optimization

OPTIMIZED QF PROTOCOL (2/5)



- **Goal:** Optimize quantization parameters (representatives v_j^R and transition levels u_i^R) to minimize the distortion between the reconstructed signal \hat{x}_{rd} and the original source signal x

$$D_{11} = \underbrace{E [(\hat{x}_{rd}^R)^2] - 2E [\hat{x}_{rd}^R x^R] + E [(x^R)^2]}_{D_{11}^R} \leftarrow \text{Real part of the signal} + \underbrace{E [(\hat{x}_{rd}^I)^2] - 2E [\hat{x}_{rd}^I x^I] + E [(x^I)^2]}_{D_{11}^I}$$

- **Uniform quantization:** Numerical optimization (Δ).

OPTIMIZED QF PROTOCOL (3/5)

○ Optimal quantization :

- Explicit expression for optimal representatives no matter the modulation types at S and R.
- Explicit expression for optimal transition levels when using simple modulations at S (ex. of 4-QAM $(x^R, x^I) \in \{-A, +A\}^2$) .

$$v_\ell^{R,*} = A \times \frac{\sum_{k=1}^L P_{k,\ell}^R \int_{u_k^R}^{u_{k+1}^R} \phi(t-A) - \phi(t+A) dt}{\sum_{k=1}^L P_{k,\ell}^R \int_{u_k^R}^{u_{k+1}^R} \phi(t-A) + \phi(t+A) dt} \quad (1)$$

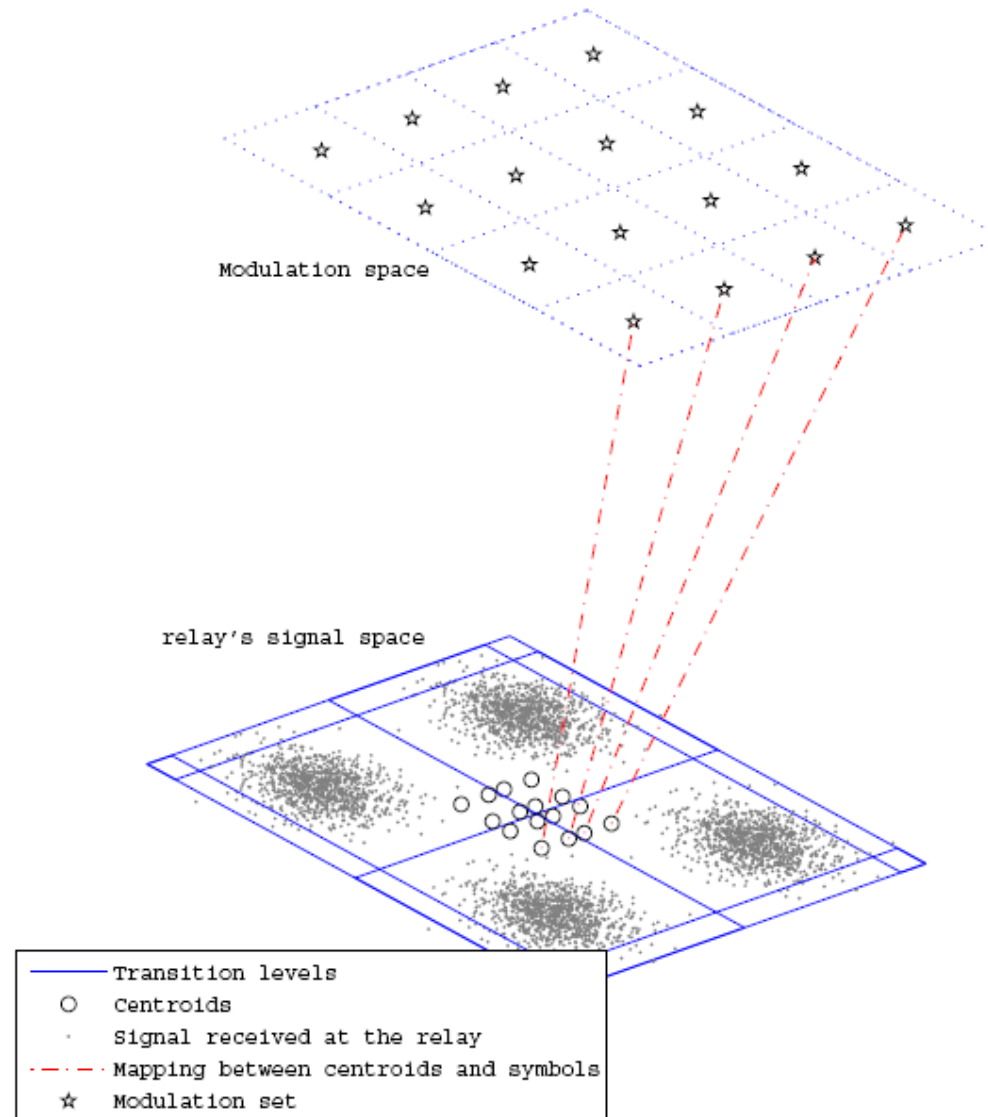
↑ **R-D link** ↑ **S-R link**

$$u_\ell^{R,*} = \frac{\sigma_{sr}^2}{2A} \ln \left[\frac{\sum_{k=1}^L (P_{\ell,k}^R - P_{\ell-1,k}^R) \left(A + \frac{1}{2} v_k^R \right) v_k^R}{\sum_{k=1}^L (P_{\ell,k}^R - P_{\ell-1,k}^R) \left(A - \frac{1}{2} v_k^R \right) v_k^R} \right] \quad (2)$$

where $\phi(t) = \frac{|h_{sr}|}{\sqrt{\pi\sigma_{sr}}} \exp\left(-\frac{|h_{sr}|^2 t^2}{\sigma_{sr}^2}\right)$ is the PDF of the noise real part on the S-R link.

OPTIMIZED QF PROTOCOL (4/5)

○ Mapping :



OPTIMIZED QF PROTOCOL (5/5)

○ Iterative optimization procedure:

- **Step 1:** Set $i = 0$. Set $\epsilon = 1$. Initialize \mathcal{V}^R and \mathcal{U}^R with the Kunterbach's algorithm ($\mathcal{V}_{(0)}^R$ and $\mathcal{U}_{(0)}^R$).
- **Step 2:** Set $i \rightarrow i+1$. Fixed partition $\mathcal{U}_{(i-1)}^R \rightarrow$ eq. (1) $\mathcal{V}_{(i)}^R$. Fixed codebook $\mathcal{V}_{(i)}^R \rightarrow$ eq. (2) $\mathcal{U}_{(i)}^R$. If the realizability condition $u_1^R \leq u_2^R \dots \leq u_L^R$ is not met stop the procedure and keep the transition levels provided by the previous iteration.
- **Step 3:** Update ϵ as follows $\epsilon = \frac{\sum_{k=1}^L |v_{k(i)}^R - v_{k(i-1)}^R|}{\sum_{k=1}^L |v_{k(i)}^R|}$. If $\epsilon \geq \epsilon_{max}$ then go to Step 2; stop otherwise.

○ Convergence analysis:

- The convergence to at least a **local minimum** is proved if one assumes a zero-mean channel input and a gaussian noise on the backward channel (S-R) under the following condition

$$\forall \ell, E[\hat{X}_{rd}^R | \hat{X}_{sr}^R = v_{\ell+1}^R] > E[\hat{X}_{rd}^R | \hat{X}_{sr}^R = v_{\ell}^R]$$

MAXIMUM LIKELIHOOD COMBINER FOR QF (1/2)

- A conventional MRC can lead to a performance loss due to quantization noise when the relay is in bad condition with a low number of quantization bits.

- The likelihood $p_{ML} = p(y_{sd}, \hat{x}_{rd}|x)$ can be factorized as:

$$p_{ML} = p(y_{sd}|x)p(\hat{x}_{rd}|x)$$

- Where $p(y_{sd}|x) = \frac{1}{\pi\sigma_{sd}^2} \exp\left(-\frac{|y_{sd}-h_{sd}x|^2}{\sigma_{sd}^2}\right)$ and

$$\begin{aligned} p(\hat{x}_{rd} = v_i|x) &= \int_{x_{sr}} p(x_{sr}, \hat{x}_{rd} = v_i|x) dx_{sr} \\ &= \int_{x_{sr}} p(x_{sr}|x) p(\hat{x}_{rd} = v_i|x_{sr}) dx_{sr} \\ &= \sum_{\ell=1}^{\sqrt{M_r}} \sum_{m=1}^{\sqrt{M_r}} P_{j,i} \left[\int_{u_\ell^R}^{u_{\ell+1}^R} \phi(t - x^R) dt \int_{u_m^I}^{u_{m+1}^I} \phi(t' - x^I) dt' \right] \end{aligned}$$

MAXIMUM LIKELIHOOD COMBINER FOR QF (2/2)

- Denote by $\underline{s} = (s_1, \dots, s_N)$ the vector of bits associated with the source symbol x .
- The log-likelihood ratio associated to the n^{th} bit:

$$\lambda(s_n) = \log \left[\frac{\sum_{\underline{s} \in S_1^{(n)}} p(y_{sd} | x) p(\hat{x}_{rd} | x)}{\sum_{\underline{s} \in S_0^{(n)}} p(y_{sd} | x) p(\hat{x}_{rd} | x)} \right]$$

where

$$S_1^{(n)} = \{(s_1, \dots, s_N) \in \{0, 1\}^N \mid s_n = 1\}$$

$$S_0^{(n)} = \{(s_1, \dots, s_N) \in \{0, 1\}^N \mid s_n = 0\}$$

If $\lambda(s_n) > 0$ then $\hat{s}_n = 1$ and $\hat{s}_n = 0$ otherwise

CLIPPED AF PROTOCOL (1/3)

- **Advantage:** Simple clipping function depending on only one parameter to be optimized β
- **Clipping function at the relay:** Real part of the signal

$$f_{\beta}^R(x^R) = \begin{cases} x^R & |x^R| \leq \beta \\ \beta \cdot \text{sgn}(x^R) & |x^R| > \beta \end{cases}$$

↓ Transmitted signal at the relay

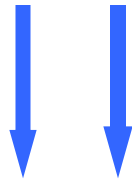
$$x_r(n) = \alpha \left[f_{\beta}^R(x_{sr}^R(n)) + j \cdot f_{\beta}^I(x_{sr}^I(n)) \right]$$

↑
Function of β : satisfy the unit-power constraint

CLIPPED AF PROTOCOL (2/3)

- Optimization criteria: Minimization of the end-to-end distortion

$$\begin{aligned} J(\beta) &= E \left[\left| \frac{1}{\alpha_\beta} \frac{h_{rd}^*}{|h_{rd}|^2} y_{rd} - x \right|^2 \right] \\ &= \frac{2}{\sqrt{M_s}} \sum_{x^R} \left\{ E \left[f_\beta^2(x^R + w_{sr}^R/h_{sr}) - 2x^R f_\beta(x^R + w_{sr}^R/h_{sr}) \right] + \frac{\sigma_{rd}^2}{\alpha_\beta^2 |h_{rd}|^2} + 1 \right\} \end{aligned}$$



Numerical optimization based on the knowledge of the CSI and SNRs at the relay for its **backward** and **forward** channels.

CLIPPED AF PROTOCOL (3/3)

○ **Let** $z \sim \mathcal{N}(\mu, \sigma^2)$ **and** $f_\beta(\cdot)$ **be the clipping function.**

For $\beta = 1$ **,** **the first and the second order moments of a clipped gaussian signal are given by:**

$$\begin{aligned} E[f_1(z)] &= \frac{1}{\sqrt{2\pi}\sigma^2} \int_{-\infty}^{\infty} f_1(z) e^{-\frac{(z-\mu)^2}{2\sigma^2}} dz \\ &= \mu + \frac{1}{\sqrt{2\pi}} \sigma e^{-\frac{(1+\mu)^2}{2\sigma^2}} - e^{-\frac{(1-\mu)^2}{2\sigma^2}} - \mu \left[Q\left(\frac{1+\mu}{\sigma}\right) + Q\left(\frac{1-\mu}{\sigma}\right) \right] - Q\left(\frac{1+\mu}{\sigma}\right) + Q\left(\frac{1-\mu}{\sigma}\right) \end{aligned}$$

$$\begin{aligned} E[f_1^2(z)] &= \frac{1}{\sqrt{2\pi}\sigma^2} \int_{-\infty}^{\infty} f_1^2(z) e^{-\frac{(z-\mu)^2}{2\sigma^2}} dz \\ &= (\sigma^2 + \mu^2) - \frac{1}{\sqrt{2\pi}} \sigma \left[(1+\mu) e^{-\frac{(1-\mu)^2}{2\sigma^2}} + (1-\mu) e^{-\frac{(1+\mu)^2}{2\sigma^2}} \right] + (1 - \sigma^2 - \mu^2) \left[Q\left(\frac{1+\mu}{\sigma}\right) + Q\left(\frac{1-\mu}{\sigma}\right) \right] \end{aligned}$$

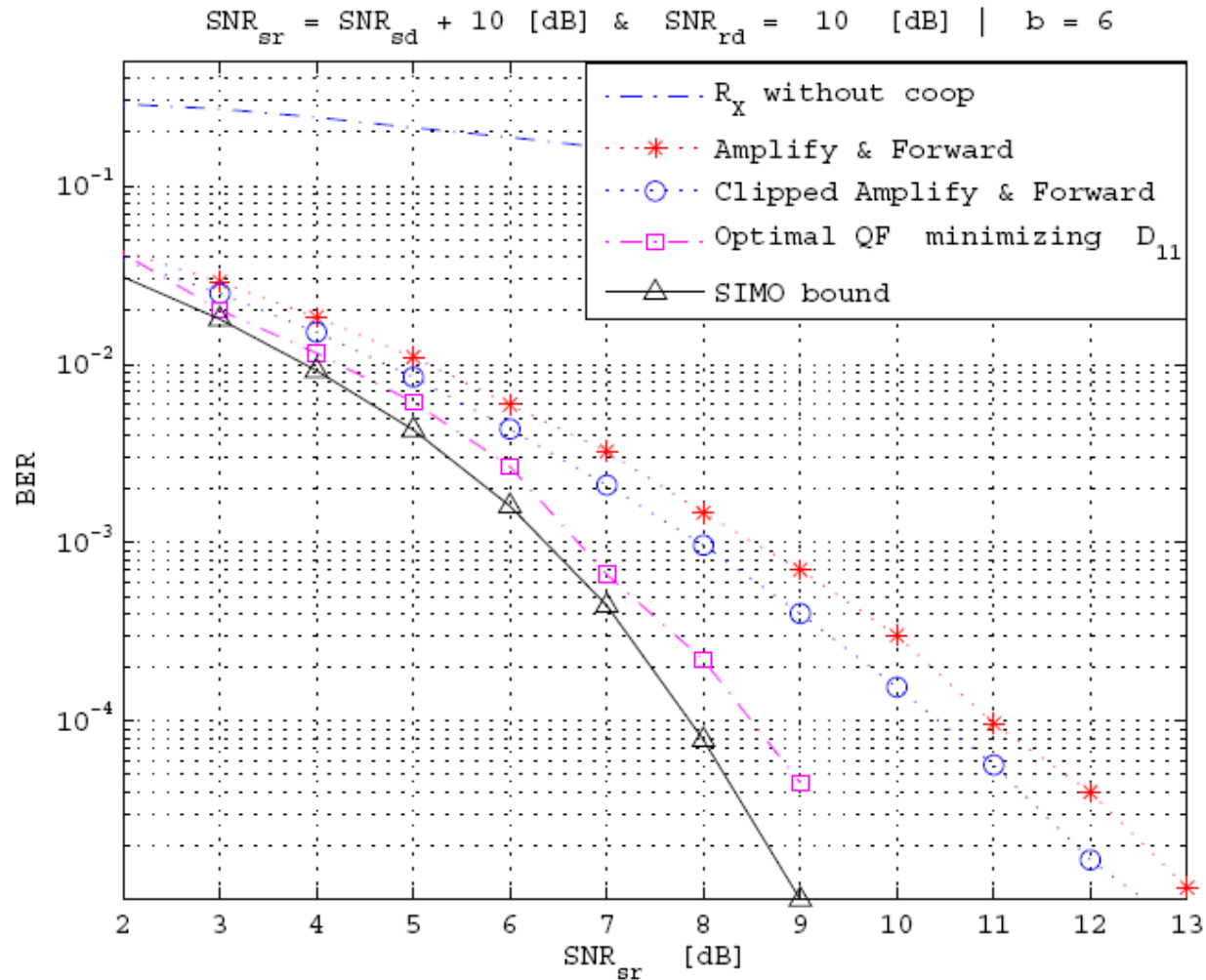
where Q **is the classical error function:** $Q(x) = \frac{1}{\sqrt{2\pi}} \int_x^{\infty} e^{-t^2} dt$

SIMULATION RESULTS (1/4)

- Simulation parameters:
 - Constant spectral efficiency.
 - We consider a 4-QAM source and static (gaussian or purely rician) or quasi-static (Rayleigh block fading model) channels.
 - Number of quantization bits for the QF: $b \in \{2, 6\} \rightarrow 2^b\text{-QAM}$ at the relay.
 - For reference we consider the case when no relay is available and the full cooperation ($\sigma_{rd} \rightarrow 0$).

SIMULATION RESULTS (2/4)

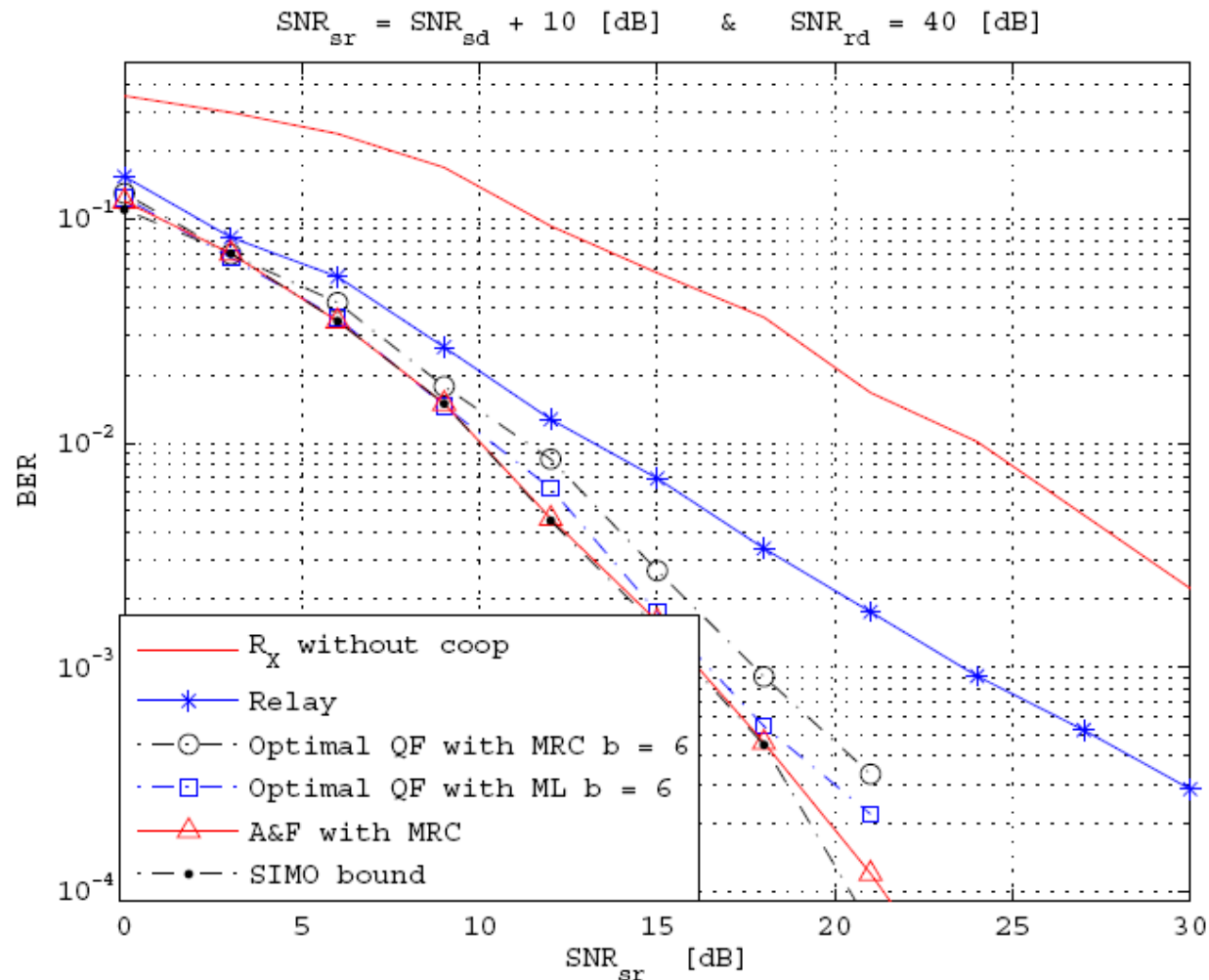
- Static channels with MRC combiner for all schemes
 - QF → Significant gain over AF and C-AF
 - Clipping gain



SIMULATION RESULTS (3/4)

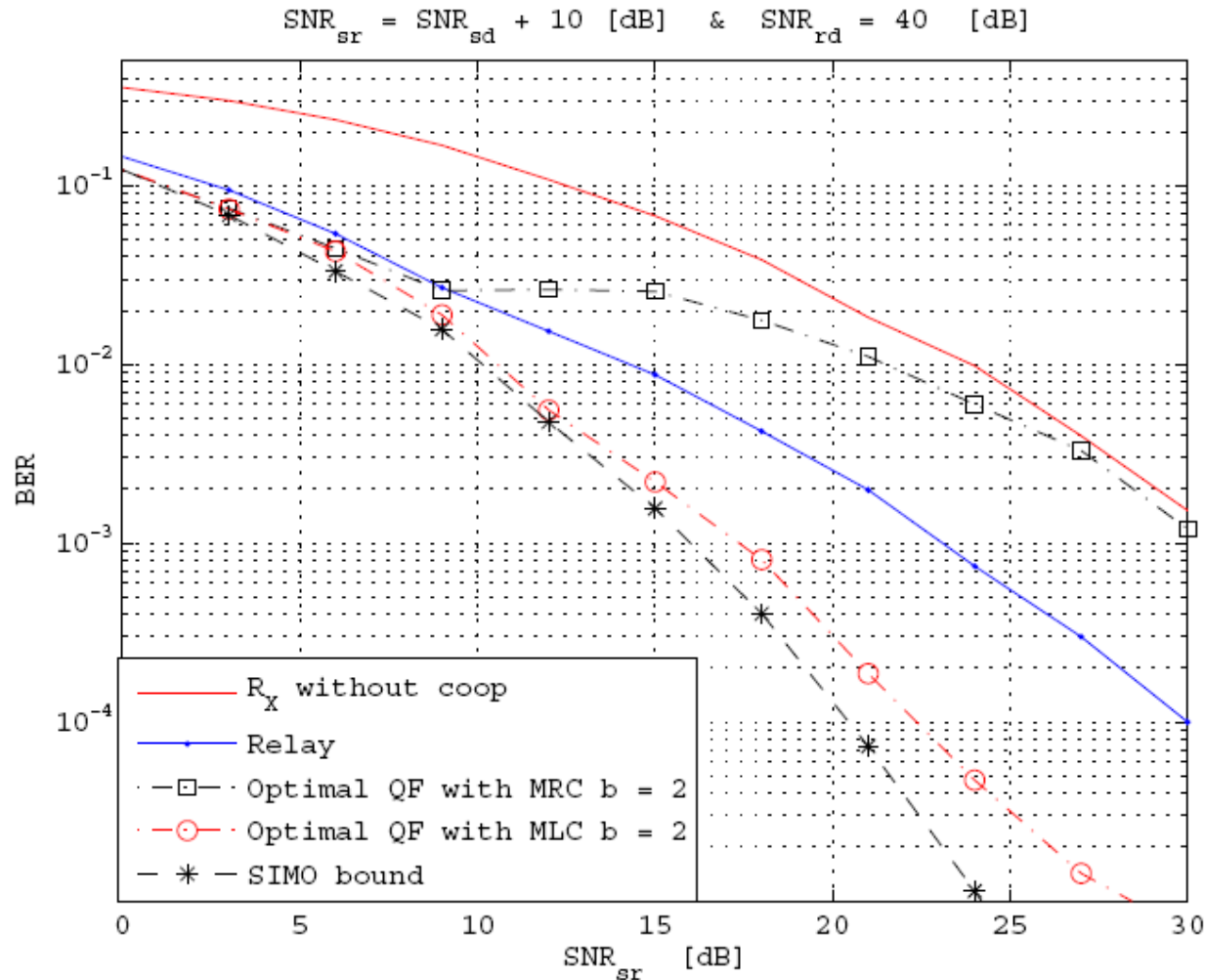
- Block fading channels

- Quasi-similar performance for these relaying schemes



SIMULATION RESULTS (4/4)

- Block fading channels: influence of the combining scheme
 - MLD → significant performance gain
 - Receiver structure adapted to the assumed relaying scheme



CONCLUSION

- **Low-complexity QF:**
 - knowledge of CSIs and SNRs at the relay.
 - Better or similar performance in static channels.
- **Clipped AF**
 - Knowledge of CSIs and SNRs at the relay.
 - Performance gain w.r.t. the conventional AF.
- **ML detector**
 - QF protocol → similar performance in quasi-static fading channel
- **QF protocol** → channel optimized AF-type protocol in a digital relay transceiver
- **Extension:** exploit the structure of channel coding.

COMPLEXITY ANALYSIS

	Optimal Quantizer	Uniform Quantizer
Creation*	$\max\{\mathcal{O}(cL^2 A^{2/3}), \mathcal{O}(cS\sqrt{M_s}LA^{2/3}), \mathcal{O}(cS\sqrt{M_s}L^2)\}$	$\max\{\mathcal{O}(SL^2\sqrt{M_s}), \mathcal{O}(S\sqrt{M_s}LA^{2/3})\}$
Storage*	$\mathcal{O}(L)$	$\mathcal{O}(L)$
Computation**	$\mathcal{O}(L)$	$\mathcal{O}(L)$

(): per SNR value; (**): per symbol to quantize*

c: nb of iterations, A: accuracy in nb of used digits, S: nb of tested points in the exhaustive search